# Lecture 4: Introduction to Cellular Automata 

## Complex Systems 530

## What is a cellular automaton?

- Automata: "a theoretical machine that changes its internal state based on inputs and its previous state" (usually finite and discrete) - Sayama p. 185
- Cellular automata: automata on a regular spatial grid, that update state based on their neighbors' states, using a state transition function
- Usually synchronous, discrete in time \& space, often deterministic (but not always!)

Neighborhood


State set


Configulation at time $t+1$


State-transition function

| CTRBL | CTRBL | CTRBL | CTRBL |
| :--- | :--- | :--- | :--- |
| $\square \square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square \square-\square$ | $\square \square \square \square-\square$ |
| $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square \square-\square$ | $\square \square \square-\square$ |
| $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square-\square$ |
| $\square \square \square \square \square$ | $\square \square \square \square \square$ | $\square \square \square \square-\square$ | $\square \square \square \square-\square$ |
| $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square-\square$ |
| $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square-\square$ | $\square \square \square-\square$ |
| $\square \square \square \square \square$ | $\square \square \square \square-\square$ | $\square \square \square \square-\square$ | $\square \square \square \square \square$ |
| $\square \square \square-\square$ | $\square \square \square \square-\square$ | $\square \square \square \square-\square$ | $\square \square \square-\square$ |

Figure 11.1: Schematic illustration of how cellular automata work.

## Cellular automata

- Cellular automata can generate highly nonlinear, even seemingly random behavior
- Much more complexity than one might expect from simple rules-emergent behavior
- To explore this, let's start with an even 'simpler’ type of cellular automata-1-dimensional CA and some of the classic work of Stephen Wolfram


## 1-dimensional CA

- We can think of our grid as a string or line of cells
- Finite sequence - 1 row of cells, so everyone has 2 neighbors except the end points
- Choose how to interpret the ends (lack of neighbors or fixed states at ends)
- Ring - all cells have 2 neighbors
- Infinite sequence - an infinite number of cells arranged in a row


## Finite sequence 1D CA

- Start with a 3-cell neighborhood (left, self, right)
- We can fully specify our CA by listing all the possible neighborhood configurations and saying what happens to the center cell, for example:

| prev | 111 | 110 | 101 | 100 | 011 | 010 | 001 | 000 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| next | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 |

- We can name our CA by translating the "next" row from binary to decimal: this is Rule 50!
(256 total possible CAs of this type)


## Rule 50



Figure 6.1: Rule 50 after 10 time steps.

## Rule 30

rule 30


What happens if we keep going?


## Wolfram's CA Classification

- CA can produce surprisingly complex behavior
- Wolfram classification - 4 classes of 1D CA
- Class I - almost all initial conditions evolve to a homogeneous state, any initial randomness is lost (e.g. Rule 0)
- Class II - Simple pattern, stable, oscillating, nested structure (e.g. Rule 18)


Figure 6.3: Rule 18 after 64 steps.

## Wolfram's CA Classification

- Class III - CAs that produce seemingly random or chaotic patterns
- Can produce sequences difficult to distinguish statistically from random, though the underlying process is deterministic
- Class III CAs typically do not


Figure 6.4: Rule 30 after 100 time steps. produce long-lasting structures (persisting over many time steps)

## Wolfram's CA Classification

- Class IV - Evolve in complex ways that involve a mix of "chaotic" and "ordered" (Class II and Class III)
- Have the potential to evolve local structures that persist over many


Figure 6.5: Rule 110 after 100 time steps. time steps



Figure 6.6: Rule 110 with random initial conditions and 600 time steps.

# Class IV CA's and computability 

- Rule 110 has been proved to be computationally universal, i.e. Turing complete (Cook M., 1998)
- So is Conway's Game of Life (classic 2D CA), and others
- Such CA can be used to compute any computable function (discuss Church-Turing Thesis)
- Wolfram's Conjecture: Every Class IV CA is Turing complete?


## Cellular Automata

- Dimensionality - How many dimensions?
- Boundaries - none (infinite domain), periodic (wrapped), cut-off (edge cells have fewer neighbors), fixed (edge cells take a fixed state)
- Grid size
- Grid type - for 2D and higher; square is typical ( \& will be our focus), but can do others!


## Cellular Automata

- State Set - binary, n-ary?
- Initial conditions - single cell active, random, etc.
- Neighborhood - queen/rook (Moore/Von Neumann), neighborhood radius
- Rules - totalistic (depends only on sum over neighborhood, e.g. majority rule), symmetric (e.g. state transition is the same up to rotation)


## CA Notation

$$
s_{t+1}(x)=F\left(s_{t}\left(x+d x_{0}\right), s_{t}\left(x+d x_{1}\right), \ldots, s_{t}\left(x+d x_{n-1}\right)\right)
$$

- $s_{t}(x)$ is the state of cell $x$ at time $t$
- $N=\left\{d x_{0}, d x_{1}, \ldots, d x_{n-1}\right\}$ is the neighborhood
- Neighborhood usually defined as cells within a given radius $r$ of $x$


## Parity Rule

$$
s_{t+1}(x)=\sum_{i=0}^{n-1} s_{t}\left(x+d x_{i}\right) \bmod k
$$

- Based on the mod $k$ sum of neighborhood values (where $k$ is the number of states)
- For binary CA, means they turn on/off based on if sum is even/odd


## Conway's Game of Life

- Possibly the most classic/well-known CA
- Large community of researchers/hobbyists, helped kick-start the field of 'artificial life'
- Produces enormous range of interesting, non-trivial behaviors
- Turing-complete


## Conway's Game of Lie

- Queen neighborhood (Moore neighborhood)
- A dead cell becomes alive if surrounded by exactly 3 live cells
- A living cell remains alive if surrounded by 2 or 3 living cells, otherwise it dies (either due to over- or underpopulation)


## Conway's Game of Life



Time $=3$

Time $=10$

Time $=50$

Time $=100$


Figure 11.6: Typical behavior of the most well-known binary CA, the Game of Life.


-


## Conway's Game of Life

- Epic collection of Conway's Game of Life patterns: https://youtu.be/C2vgICfQawE?t=70
- Nicky Case Simulator version: https://ncase.me/sim/?s=conway
- Web version to try: https://playgameoflife.com
- ca-gameoflife.py in PyCX
- Game of life wiki: https://conwaylife.com/wiki/Main Page
- NYT: https://www.nytimes.com/2020/12/28/science/math-conway-game-of-life.html


## Turmites

- 2D Turing machine generalizations
- Named "Turmites" after Turing and the fact that the write-head of the 'machine' moves similarly to a bug
- The 'turmite' or 'ant'
- E.g. Langton's Ant


# Applications of CA \& real-world examples 

- Forest fire models/disease epidemics
- Sand heaps/avalanches
- Majority rule and voter models
- Diffusion-limited aggregation (DLA), percolation, lattice models of materials
- And many more-some more realistic than others
- Many ABMs can be viewed as CA, or near-CA (e.g. if we allow probabilistic rather than deterministic rules)


## CA on seashells

- Conus textile appears to operate with Rule 30 (or close to it)



## CA on lizard scales




## CA \& ABM Dynamics

- Not always easy to interpret! Can have many patterns, as we saw with Game of Life, etc.
- However, sometimes there are major overall patterns that we can see
- More on this next time!


## For next time...

- Reading
- Sayama Chapter 11
- Think Complexity Chapter 6
- We'll discuss 2D CA, how to build CA, variations on CA, and theory for how to analyze the complexity and dynamics of CA

